

**Quantum Computing**  
**“Uygulama”**

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# Örnek-1:

- $\psi = \begin{pmatrix} \alpha \\ \beta \end{pmatrix} = \alpha|0\rangle + \beta|1\rangle$  ise  $|0\rangle = \begin{pmatrix} 1 \\ 0 \end{pmatrix}$  ve  $|1\rangle = \begin{pmatrix} 0 \\ 1 \end{pmatrix}$  vektörleri ile temsil edilir.
- $ZR_\varphi\psi = \alpha|0\rangle - \beta|1\rangle$  olması için  $\varphi=?$
- $\begin{bmatrix} 1 & 0 \\ 0 & -1 \end{bmatrix} \begin{bmatrix} 1 & 0 \\ 0 & e^{i\varphi} \end{bmatrix} \begin{pmatrix} \alpha \\ \beta \end{pmatrix} = \begin{bmatrix} 1 & 0 \\ 0 & -1 \end{bmatrix} \begin{pmatrix} \alpha \\ \beta e^{i\varphi} \end{pmatrix} = \begin{pmatrix} \alpha \\ -\beta e^{i\varphi} \end{pmatrix} = \alpha|0\rangle - \beta e^{i\varphi}|1\rangle$
- $\varphi=0$

# Örnek-2:

- $\psi = \begin{pmatrix} \alpha \\ \beta \end{pmatrix} = \alpha|0\rangle + \beta|1\rangle$  ise  $|0\rangle = \begin{pmatrix} 1 \\ 0 \end{pmatrix}$  ve  $|1\rangle = \begin{pmatrix} 0 \\ 1 \end{pmatrix}$  vektörleri ile temsil edilir.
- $XR_\varphi\psi = \beta|0\rangle + \alpha|1\rangle$  olması için  $\varphi=?$
- $\begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} \begin{bmatrix} 1 & 0 \\ 0 & e^{i\varphi} \end{bmatrix} \begin{pmatrix} \alpha \\ \beta \end{pmatrix} = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} \begin{pmatrix} \alpha \\ \beta e^{i\varphi} \end{pmatrix} = \begin{pmatrix} \beta e^{i\varphi} \\ \alpha \end{pmatrix} = \beta e^{i\varphi}|0\rangle + \alpha|1\rangle$
- $\varphi=0$

# Örnek-3:

- $XS|1\rangle=?$

- $\begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} \begin{bmatrix} 1 & 0 \\ 0 & i \end{bmatrix} \begin{pmatrix} 0 \\ 1 \end{pmatrix} = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} \begin{pmatrix} 0 \\ i \end{pmatrix} = \begin{pmatrix} i \\ 0 \end{pmatrix} = i \begin{pmatrix} 1 \\ 0 \end{pmatrix} = i|0\rangle$

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- $XS|0\rangle=?$

- $\begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} \begin{bmatrix} 1 & 0 \\ 0 & i \end{bmatrix} \begin{pmatrix} 1 \\ 0 \end{pmatrix} = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} \begin{pmatrix} 1 \\ 0 \end{pmatrix} = \begin{pmatrix} 0 \\ 1 \end{pmatrix} = |1\rangle$

# Örnek-4:

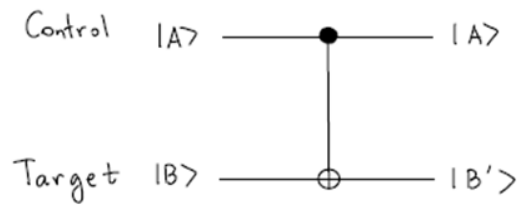
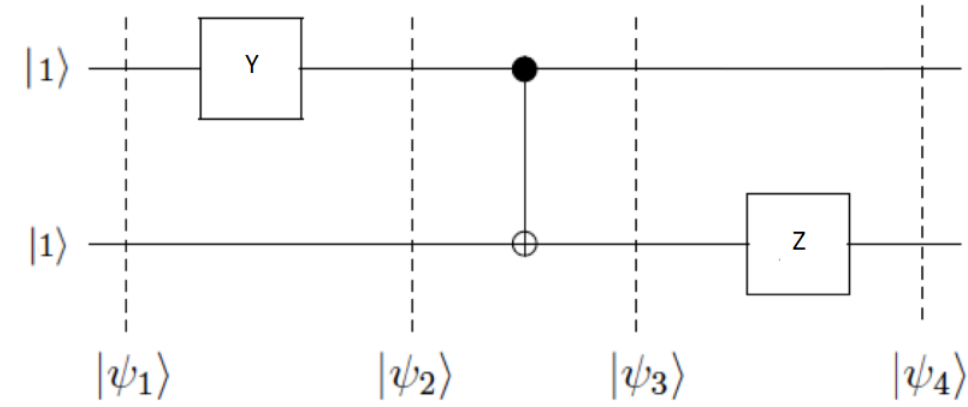
- $\begin{pmatrix} -i \\ 0 \end{pmatrix} \otimes \begin{pmatrix} 0 \\ 1 \end{pmatrix} = ?$

$$\begin{pmatrix} 0 \\ -i \\ 0 \\ 0 \end{pmatrix} = -i \begin{pmatrix} 0 \\ 1 \\ 0 \\ 0 \end{pmatrix} = -i|01\rangle$$

- $\begin{pmatrix} 0 \\ i \end{pmatrix} \otimes \begin{pmatrix} 1 \\ 0 \end{pmatrix} = ?$

$$\begin{pmatrix} 0 \\ 0 \\ i \\ 0 \end{pmatrix} = i \begin{pmatrix} 0 \\ 0 \\ 1 \\ 0 \end{pmatrix} = i|10\rangle$$

# Örnek-5:



$ AB\rangle$	$ AB'\rangle$
100	100
101	101
110	111
111	110

A kontrol girişi olduğu gibi geçer.  
B çıkışı, A=0 ise B olduğu gibi geçer, A=1 ise B'nin tersi geçer,

Üst hatta,

$$|\psi_1\rangle = |1\rangle$$

$$|\psi_2\rangle = \begin{bmatrix} 0 & -i \\ i & 0 \end{bmatrix} \begin{pmatrix} 0 \\ 1 \end{pmatrix} = \begin{pmatrix} -i \\ 0 \end{pmatrix} = -i|0\rangle$$

$$|\psi_3\rangle = -i|0\rangle$$

$$|\psi_4\rangle = -i|0\rangle$$

Alt hatta,

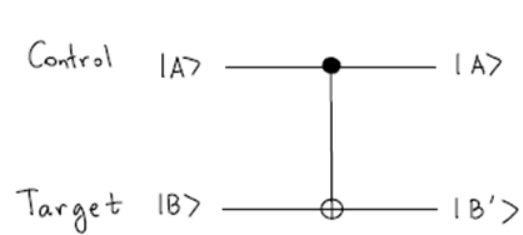
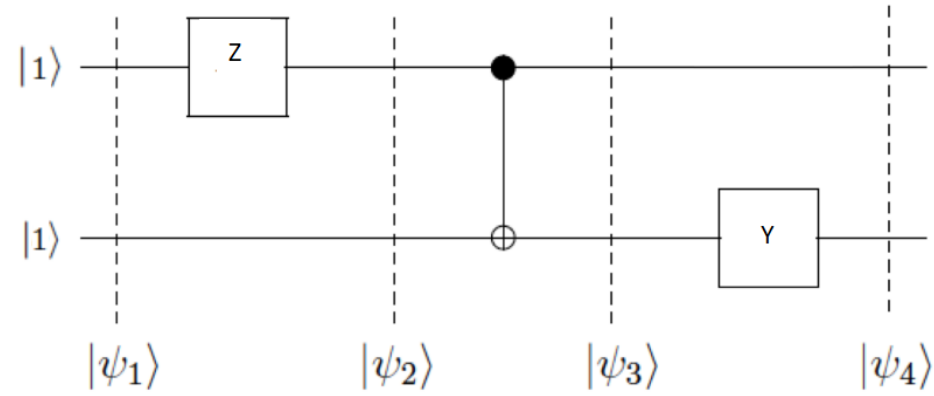
$$|\psi_1\rangle = |1\rangle$$

$$|\psi_2\rangle = |1\rangle$$

$$|\psi_3\rangle = |1\rangle$$

$$|\psi_4\rangle = \begin{bmatrix} 1 & 0 \\ 0 & -1 \end{bmatrix} \begin{pmatrix} 0 \\ 1 \end{pmatrix} = -|1\rangle$$

# Örnek-6:



$ AB\rangle$	$ AB'\rangle$
$ 00\rangle$	$ 00\rangle$
$ 01\rangle$	$ 01\rangle$
$ 10\rangle$	$ 11\rangle$
$ 11\rangle$	$ 10\rangle$

A kontrol grişi olduğu gibi geçer.  
B çıkışı, A=0 ise B olduğu gibi geçer, A=1 ise B'nin tersi geçer,

Üst hatta,

$$|\psi_1\rangle = |1\rangle$$

$$|\psi_2\rangle = \begin{bmatrix} 1 & 0 \\ 0 & -1 \end{bmatrix} \begin{pmatrix} 0 \\ 1 \end{pmatrix} = -|1\rangle$$

$$|\psi_3\rangle = -|1\rangle$$

$$|\psi_4\rangle = -|1\rangle$$

Alt hatta,

$$|\psi_1\rangle = |1\rangle$$

$$|\psi_2\rangle = |1\rangle$$

$|\psi_3\rangle = |0\rangle$ ; üst hat 1 olduğundan alt hattın tersi geçer.

$$|\psi_4\rangle = \begin{bmatrix} 0 & -i \\ i & 0 \end{bmatrix} \begin{pmatrix} 1 \\ 0 \end{pmatrix} = \begin{pmatrix} 0 \\ i \end{pmatrix} = i \begin{pmatrix} 0 \\ 1 \end{pmatrix} = i|1\rangle$$

# Örnek-6:

$$|\psi_2\rangle = \begin{pmatrix} -i \\ 0 \end{pmatrix} \otimes \begin{pmatrix} 0 \\ 1 \end{pmatrix} = \begin{pmatrix} 0 \\ -i \\ 0 \\ 0 \end{pmatrix} = -i \begin{pmatrix} 0 \\ 1 \\ 0 \\ 0 \end{pmatrix} = -i|01\rangle$$

$$|\psi_2\rangle = \begin{pmatrix} -i \\ 0 \end{pmatrix} \otimes \begin{pmatrix} 0 \\ i \end{pmatrix} = \begin{pmatrix} 0 \\ 1 \\ 0 \\ 0 \end{pmatrix} = \begin{pmatrix} 0 \\ 1 \\ 0 \\ 0 \end{pmatrix} = |01\rangle$$



# Örnek-7:

## Entangled states

As we shall see, quantum computing draws its advantage from the fact that not all quantum states are separable. Consider the *two qubit unitary*

$$\text{CNOT} = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{bmatrix}$$

applied to the state

$$|+\rangle \otimes |0\rangle = (1/\sqrt{2}) [1 \ 1]^T \otimes [1 \ 0]^T = [1/\sqrt{2} \ 0 \ 1/\sqrt{2} \ 0]^T:$$

$$\begin{aligned} \text{CNOT}(|+\rangle \otimes |0\rangle) &= \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{bmatrix} \begin{bmatrix} 1/\sqrt{2} \\ 0 \\ 1/\sqrt{2} \\ 0 \end{bmatrix} = \begin{bmatrix} 1/\sqrt{2} \\ 0 \\ 0 \\ 1/\sqrt{2} \end{bmatrix} \\ &= \frac{1}{\sqrt{2}}(|00\rangle + |11\rangle) \end{aligned}$$

This is an *an entangled state which cannot be separated as tensor product*. We call **CNOT** an *entangling* operation (or “gate” in the quantum circuit model), and even though it operates on two qubits, Postulate 2 still applies – so it is still necessarily unitary.

# Örnek-8:

## More on entangled states and the Bell states

In quantum computing, we take **non-separability as the definition of an entangled state**, and so there are infinitely many entangled states. **Four important two-qubit entangled states are the Bell states:**

$$|\Phi^+\rangle = (1/\sqrt{2})(|00\rangle + |11\rangle) = (1/\sqrt{2}) [1 \ 0 \ 0 \ 1]$$

$$|\Phi^-\rangle = (1/\sqrt{2})(|00\rangle - |11\rangle) = (1/\sqrt{2}) [1 \ 0 \ 0 \ -1]$$

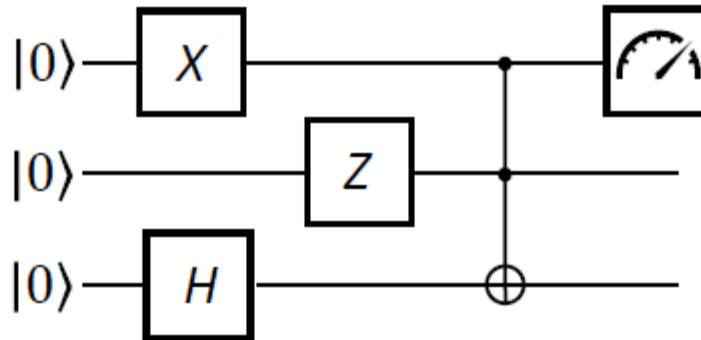
$$|\Psi^+\rangle = (1/\sqrt{2})(|01\rangle + |10\rangle) = (1/\sqrt{2}) [0 \ 1 \ 1 \ 0]$$

$$|\Psi^-\rangle = (1/\sqrt{2})(|01\rangle - |10\rangle) = (1/\sqrt{2}) [0 \ 1 \ -1 \ 0]$$

Which form an orthonormal basis for  $\mathbb{C}^4$  (**exercise: verify this**).

# Örnek-9:

## Quantum circuits



A quantum circuit is a tensor network of  $n$  qubits, with three stages:

- Initialisation of all qubits in the  $|0\rangle$  state (denoted  $|0\rangle^{\otimes n}$ ).
  - Sometimes it will be convenient to let the initial state be something other than  $|0\rangle^{\otimes n}$  – but we should be able to efficiently prepare this initial state from  $|0\rangle^{\otimes n}$  (that is, using a number of one- and two-qubit operations that is at most polynomial in the number of qubits).
- Some quantum gates, which represent unitary transformations.
- A final layer of measurements in the computational basis, on some or all of the qubits.
  - In fact, by the *principle of implicit measurement*, we can consider **all** qubits to be measured in the final layer.

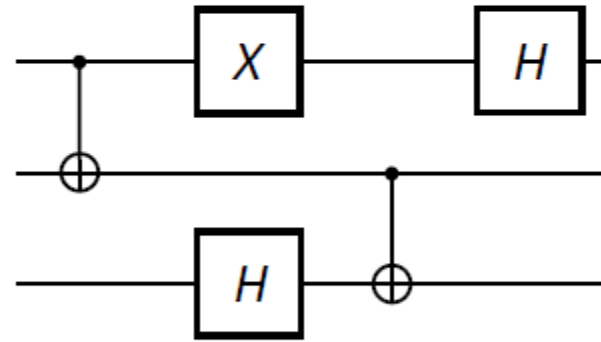
# Örnek-10:

## The matrix of a quantum circuit

As the quantum circuit (with the initialisation and measurement stages omitted) just represents a unitary evolution, **we can express the whole thing as a matrix**. We must follow the following two rules:

- Composition across wires is achieved by the tensor product.
- Composition along (sets of) wires is achieved by the normal matrix product, but **right to left**.

For example:



Is equal to:

$$(H \otimes I_4) \times (I_2 \otimes \text{CNOT}) \times (X \otimes I_2 \otimes H) \times (\text{CNOT} \otimes I_2)$$

where  $I_2$  is the  $2 \times 2$  identity, and  $I_4 = I_2 \otimes I_2$  is the  $4 \times 4$  identity.

# Örnek-11:

## A universal gate-set

Perhaps surprisingly, **only three gates are needed to form a universal gate-set**, two we have met: **CNOT** and  **$H$** , and the third is:

$$T = \begin{bmatrix} 1 & 0 \\ 0 & e^{i\pi/4} \end{bmatrix}$$

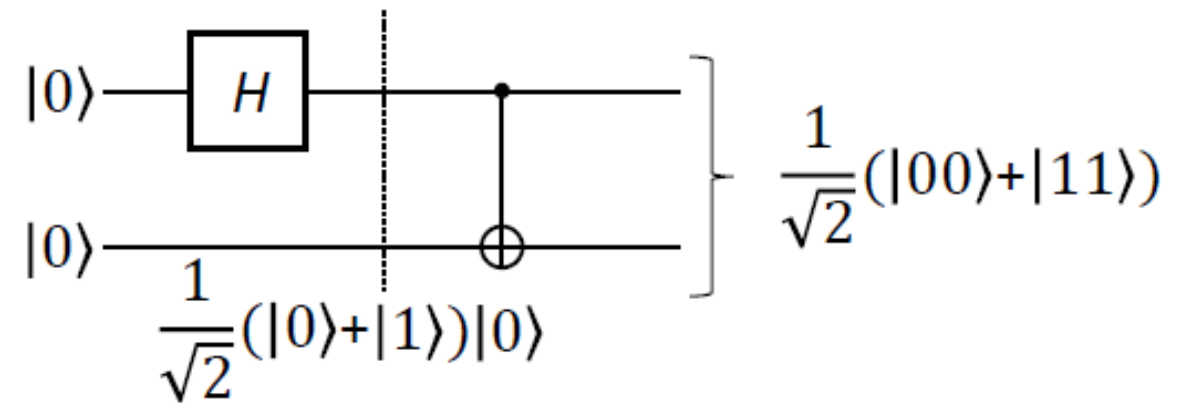
The introduction of this  $T$  gate is, however, crucial, and the famous Gottesman-Knill theorem holds that any circuit consisting of just the gates we have met thus far  **$X, Y, Z, H, S, \text{CNOT}$**  can be efficiently simulated on a classical computer.

We can see that the single-qubit gates we have met so far can be expressed in terms of  $H$  and  $T$  as follows:

- $S = T^2$
- $Z = S^2$
- $X = HZH$
- $Y = iXZ = SXSZ$

# Örnek-12:

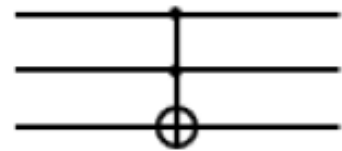
Quantum circuit example 1: entangling two qubits



# Örnek-13:

## The Toffoli gate

The Toffoli gate *does* provide a quantum generalisation of the classical **AND** gate, with three inputs and outputs.

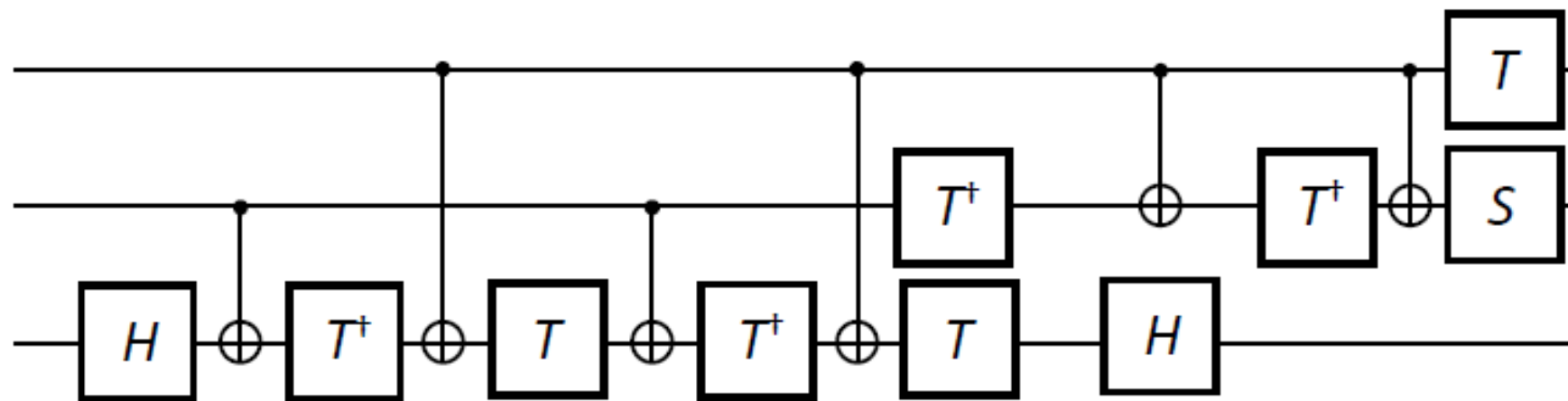


$$\begin{bmatrix} 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 \end{bmatrix}$$

When the first two inputs are classical bits ( $|0\rangle$  or  $|1\rangle$ ), and the third is  $|0\rangle$  the third output is the **AND** of the first two inputs.

# Örnek-14:

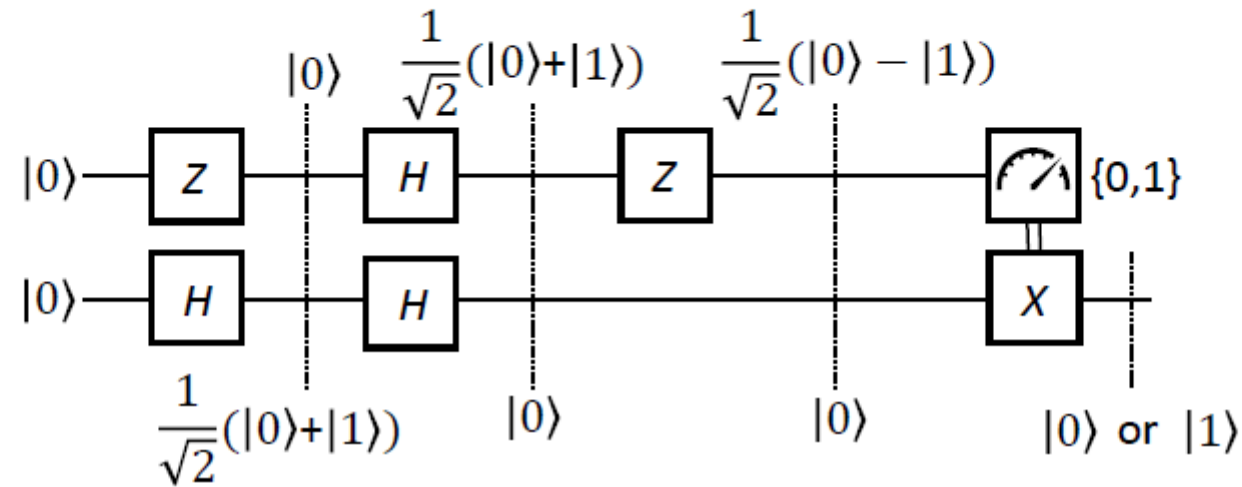
Quantum circuit example 2: decomposing the Toffoli gate into two-qubit unitaries





# Örnek-15:

Quantum circuit example 3: self-inverse nature of  $H$  and classical control



- By convention, **classical information (that is, a *bit* rather than a *qubit*) is represented by a double-line**. So we should read this classical control as: if the measurement outcome is zero, then do nothing; if the measurement outcome is one, then perform the Pauli- $X$  operation as shown.
- The inclusion of classical control may appear to contradict our definition of a quantum circuit as having measurement only in a final layer (which contains nothing other than measurement). However, the *principle of deferred measurement* means that any quantum circuit *can* be expressed in this form. But in practise it is often convenient to allow measurements to occur mid-circuit.